# Point-free geometry and topology Part VI: Point-free geometry and verisimilitude of theories

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### Outline

Pseudo-metric and metric spaces

Approximate point-free (pseudo-)metric spaces

Theories and truth

Point-free spaces of theories

Spaces of theories

Verisimilitude

### Pseudo-metric spaces

### Definition (of a pseudo-metric space)

A pseudo-metric space (pm-space for short) is a pair  $\langle X, \delta \rangle$  such that  $X \neq \emptyset$  and  $\delta$  is a function from  $X \times X$  to  $[0, +\infty)$  with the following properties:

The function  $\delta$  will be called pseudo-metric (p-metric), the number  $\delta(x, y)$  pseudo-distance (p-distance) between x and y.

### Examples of pm-spaces

### Example (The trivial pm-space) A pair $\langle X, \delta \rangle$ such that for all $x, y \in X$ , $\delta(x, y) = 0$ .

Example (The absolute value pm-space) A pair  $\langle \mathbb{R}^2, \delta \rangle$  with  $\delta \colon \mathbb{R}^2 \times \mathbb{R}^2 \longrightarrow [0, +\infty)$  such that:

$$\delta(\langle x_1, y_1 \rangle, \langle x_2, y_2 \rangle) \coloneqq |x_1 - x_2|.$$



### Approximate pseudo-metric spaces

Definition (of a point-free pseudometric space)

A point-free aproximate pseudometric space is a structure  $\Re = \langle \mathbf{R}, \sqsubseteq, \delta, \Delta \rangle$  such that:

$$\langle \mathbf{R}, \sqsubseteq \rangle$$
 is a poset (APM0)

and  $\delta: \mathbf{R} \times \mathbf{R} \longrightarrow [0, +\infty)$  and  $\Delta: \mathbf{R} \longrightarrow [0, +\infty)$  are such that (set  $\operatorname{diam}(x) \coloneqq \Delta(x, x)$ ):

$$\delta(x,x) = 0, \qquad (APM1)$$

$$\delta(x, y) = \delta(y, x),$$
 (APM2)

$$\Delta(x,y) = \Delta(y,x)$$
 (APM3)

$$0 \leq \Delta(x, y) - \delta(x, y) \leq \operatorname{diam}(x) + \operatorname{diam}(y), \qquad (\mathsf{APM4})$$
  
$$\delta(x, y) \leq \delta(x, z) + \delta(z, y) + \operatorname{diam}(z), \qquad (\mathsf{APM5})$$

$$\Delta(x,y) \leq \Delta(x,z) + \Delta(z,y),$$
 (APM6)

- $x_1 \sqsubseteq x_2 \land y_1 \sqsubseteq y_2 \Longrightarrow \delta(x_2, y_2) \leqslant \delta(x_1, y_1), \qquad (\mathsf{APM7})$
- $x_1 \sqsubseteq x_2 \land y_1 \sqsubseteq y_2 \Longrightarrow \Delta(x_1, y_1) \le \Delta(x_2, y_2).$  (APM8)

### Approximate metric spaces

#### Definition (of an approximate metric space)

An apm-space  $\Re = \langle \mathbf{R}, \sqsubseteq, \delta, \Delta \rangle$  is an approximate metric space (abbr: am-space) in case it satisfies the following condition:

$$\Delta(x,y) = 0 \Longrightarrow x = y. \tag{AM}$$

### Definition (of a point-region)

Let  $\Re = \langle \mathbf{R}, \sqsubseteq, \delta, \Delta \rangle$  be an apm-space. A region  $x \in \mathbf{R}$  is a point-region iff diam(x) = 0. The set  $PR(\mathbf{R})$  is the set of all point-regions of  $\Re$ :

$$PR(\mathbf{R}) \coloneqq \{x \in \mathbf{R} \mid diam(x) = 0\}. \qquad (df PR(\mathbf{R}))$$

### Distances between regions



# **Diameters of regions**



# Theories

### Definition (of a theory)

A theory is a set of sentences.

Definition (of a true theory)

A theory is true iff all its sentences are true.

#### Definition (of a false theory)

A theory is false iff at least one of its sentences is false.

### Definition (of a consequence of a theory)

A sentence p is a consequence of a theory T iff it follows from T (whenever theory T is true, then p must be true).

#### Fact

- ► If a theory T is true, then all its consequences are true.
- If p is a consequence of T and p is false, then T is false as well.

### Definition

A theory  $T_2$  is closer to the truth than a theory  $T_1$  iff

- ▶ all true consequences of T₁ are also consequences of T₂
- ▶ all false consequences of T₂ are also consequences of T₁
- either there is some true consequence of  $T_2$  which is not a consequence of  $T_1$  or there is some false consequence of  $T_1$  which is not a consequence of  $T_2$ .

# First-order theories

- Let  $\mathcal{L}$  be any first-order language.
- Let  $For(\mathcal{L})$  be the set of all formulas of  $\mathcal{L}$ .

Definition (of a first-order theory) A first-order theory built in language  $\mathcal{L}$  is a triple:

 $T\coloneqq \langle \mathcal{L}, X, {\bf F} \rangle$ 

such that

- $X \subseteq For(\mathcal{L})$ , it will be called the set of specific axioms of T
- ► ⊢ is a provability (derivability) relation.

Remark

For a fixed theory T the set of its specific axioms will be denoted by Axm(T).

# Some special sets of formulas

Definition

Let  $X \subseteq For(\mathcal{L})$  for some first-order language  $\mathcal{L}$ . We define:

the set of formulas provable from X:

 $\operatorname{Prov}(\mathbf{X}) \coloneqq \{ F \in \operatorname{For}(\mathcal{L}) \mid \mathbf{X} \vdash F \}.$ 

the set of formulas refutable from X:

 $\operatorname{Ref}(\mathbf{X}) \coloneqq \{ F \in \operatorname{For}(\mathcal{L}) \mid \mathbf{X} \vdash \neg F \}.$ 

the set of formulas which are decidable on base of X:

 $Dec(X) \coloneqq Prov(X) \cup Ref(X)$ .

► the elements of CDec(X) will be called undecidable formulas given X.

### Definition (of a consistent set of formulas)

A set  $X \subseteq For(\mathcal{L})$  is consistent iff no formula is both provable and refutable from X:

$$\operatorname{cons}(X) \stackrel{\mathrm{df}}{\longleftrightarrow} \operatorname{Prov}(X) \cap \operatorname{Ref}(X) = \emptyset.$$
 (df  $\operatorname{cons}(X)$ )

# Binary operations on sets of formulas

### Definition

Let  $X,Y\subseteq For(\mathcal{L})$  for some first-order language  $\mathcal{L}.$  We define:

► the formulas on which X and Y agree:

 $Agree(X, Y) \coloneqq (Prov(X) \cap Prov(Y)) \cup (Ref(X) \cap Ref(Y)).$ 

▶ the formulas on which X and Y disagree:

 $Disagree(X, Y) \coloneqq (Prov(X) \cap Ref(Y)) \cup (Prov(Y) \cap Ref(X)).$ 

We also set the following terminology:

- CAgree(X, Y) will be called the set on which X and Y do not agree or may disagree,
- CDisagree(X, Y) will be called the set on which X and Y do not disagree or may agree.

## Some important facts

If X and Y are both consistent, then:

Agree $(X, Y) \cap Disagree(X, Y) = \emptyset$ , (1)

$$Agree(X, Y) \subseteq CDisagree(X, Y), \qquad (2)$$

$$Disagree(X, Y) \subseteq CAgree(X, Y),$$
(3)

$$Dec(X) \cap Dec(Y) = Agree(X, Y) \cup Disagree(X, Y),$$
 (4)

$$Disagree(X, Y) = CAgree(X, Y) \cap Dec(X) \cap Dec(Y),$$
 (5)

$$Disagree(X, Y) \subseteq Dec(X) \cap Dec(Y)$$
, (6)

 $Disagree(X, Y) \cap (CDec(X) \cup CDec(Y)) = \emptyset.$  (7)

### Tests and measures

### Definition (of a sentence)

 $F \in For(\mathcal{L})$  is a sentence iff *F* has no free variables. Let  $Sent(\mathcal{L})$  be the set of all sentences built in  $\mathcal{L}$ .

#### Definition (of a test)

Let  $F \in \text{Sent}(\mathcal{L})$ . By a test we mean the set of all formulas which are equivalent either to F or to its negation:

```
\|F\| \coloneqq \left\{ G \in \operatorname{Sent}(\mathcal{L}) \mid T \vdash G \equiv F \text{ or } T \vdash G \equiv \neg F \right\}. \qquad (df \| \|)
```

For a set  $X \subseteq Sent(\mathcal{L})$  let:

 $||\mathbf{X}|| \coloneqq \left\{ ||F|| \mid F \in \mathbf{X} \right\}.$ 

Tests and measures

#### Definition (of a vacuous test)

Let  $\perp$  be an arbitrary (but fixed) logically false formula. By a vacuous test we mean the set  $\|\perp\|$ :

```
\mathbf{0} \coloneqq \|\bot\| \ . \tag{df 0}
```

### Definition By a relevance measure we mean a function

```
rel: ||Sent(\mathcal{L})|| \longrightarrow [0, 1]
```

such that:

 $\operatorname{rel}(\mathbf{0}) \coloneqq \mathbf{0},$  $\sum_{x \in ||\operatorname{Sent}(\mathcal{L})||} \operatorname{rel}(x) \coloneqq \mathbf{1}.$ 

(df rel)

### Tests and measures

Definition (of a relevant test, relevant sentence and sensitive function)

- $X \in ||\text{Sent}(\mathcal{L})||$  is relevant iff  $\operatorname{rel}(X) \neq 0$ .
- A sentence F is relevant iff ||F|| is relevant.
- A function rel is sensitive iff all non-vacuous tests are relevant.

Definition (of a measure of the degree of relevance) By a measure of the degree of relevance of  $X \subseteq ||Sent(\mathcal{L})||$  (with respect to some rel function) we mean the function:

 $\mu \colon \mathcal{P}(\|\operatorname{Sent}(\mathcal{L})\|) \longrightarrow [0,1]$ 

such that:

$$\mu(X) \coloneqq \sum_{x \in \mathbf{X}} \operatorname{rel}(x) \,. \tag{df}\, \mu)$$

# Distances between sets of sentences

### Definition (of distances between sets of sentences) Let rel be some relevance measure.

Let

 $d\colon \mathcal{P}(\operatorname{Sent}(\mathcal{L})) \times \mathcal{P}(\operatorname{Sent}(\mathcal{L})) \longrightarrow [0,1]$ 

be such that:

$$d(X_1, X_2) \coloneqq \mu \left( \| \text{Disagree}(X_1, X_2) \| \right).$$
(8)

On the other hand, let

 $D: \mathcal{P}(\operatorname{Sent}(\mathcal{L})) \times \mathcal{P}(\operatorname{Sent}(\mathcal{L})) \longrightarrow [0, 1]$ 

be such that:

$$D(X_1, X_2) \coloneqq \mu \left( \| \big( Agree(X_1, X_2) \| \big). \right).$$
(9)

### Definition (of the actual and possible contrasts)

For given  $X_1, X_2 \subseteq Sent(\mathcal{L})$ :

- the number d(X<sub>1</sub>, X<sub>2</sub>) is called the actual contrast between the two sets in question,
- ► the number D(X<sub>1</sub>, X<sub>2</sub>) is called the possible contrast between them.

### Distances between sets of sentences

#### Definition (of the diameter of the set of sentences)

Let  $X \subseteq \text{Sent}(\mathcal{L})$ . According to the earlier definitions of the diameter of a given set and of the possible contrast between sets of sentences we have:

diam(X) :=  $D(X, X) := \mu(||CAgree(X, X)||) = \mu(||CDec(X)||)$ .

#### Fact

From the definition of  $\mu$  it follows that:

diam(X) = 
$$1 - \mu(||\text{Dec}(X)||)$$
. (10)

## Spaces of sets sentences

#### Definition

Let **CA** be the set of all consistent sets of sentences built in first-order language  $\mathcal{L}$ :

$$CA := \{X \in Sent(\mathcal{L}) \mid cons(X)\}. \qquad (df CA)$$

Definition Let  $\langle CA, \sqsubseteq \rangle$  be such that:

$$X \sqsubseteq Y \stackrel{df}{\Longleftrightarrow} Y \subseteq X \,.$$

#### Fact

The structure  $(CA, \sqsubseteq, d, D)$  is an apm-space.

Fact

If rel is sensitive, then:

$$\mu(||X||) = 1 \iff \operatorname{Rel}(\mathcal{L}) \subseteq X.$$

#### Theorem

The following equivalences hold in  $(CA, \sqsubseteq, d, D)$ :

- 1.  $d(X, Y) = 1 \iff \text{Disagree}(X, Y) = \text{Rel}(\mathcal{L}),$
- 2.  $d(X, Y) = 0 \iff \text{Disagree}(X, Y) = \emptyset$ ,
- 3.  $D(X,Y) = 1 \iff \operatorname{Agree}(X,Y) \cap \operatorname{Rel}(\mathcal{L}) = \emptyset$ ,
- 4.  $D(X,Y) = 0 \iff \operatorname{Agree}(X,Y) = \operatorname{Sent}(\mathcal{L}),$
- 5. diam $(X) = 1 \iff \text{Dec}(X) \cap \text{Rel}(\mathcal{L}) = \emptyset$ ,
- 6. diam $(X) = 0 \iff Dec(X) = Sent(\mathcal{L}).$

### Definition (of the set of theorems)

Suppose T is a first-order theory. The set of theorems of T is the set of all formulas which are provable from the specific axioms of T:

 $\mathrm{Th}(\mathrm{T}) \coloneqq \{F \in \mathrm{For}(\mathcal{L}) \mid \mathrm{Axm}(\mathrm{T}) \vdash F\}. \qquad (\mathrm{df} \, \mathrm{Th}(\mathrm{T}))$ 

That is Th(T) = Prov(Axm(T)).

### Definition (of the set of deductive closures)

Let **Th** be the set of all closures of sets of sentences built in a first-order language  $\mathcal{L}$ :

 $\mathbf{Th} := \{ Th(T) \mid T \text{ is a first-order theory} \}. \qquad (df \mathbf{Th})$ 

Of course,  $\mathbf{Th} = \{ \operatorname{Prov}(X) \mid X \subseteq \operatorname{Sent}(\mathcal{L}) \}.$ 

#### Remark

From now on by a theory we mean any element of **Th**, that is any set of theorems of some first-order theory built in a language  $\mathcal{L}$ .

#### Definition (of the apm-space of all theories)

The subspace  $\langle \mathbf{Th}, \sqsubseteq, d, D \rangle$  of the space  $\langle \mathbf{CA}, \sqsubseteq, d, D \rangle$  will be called the apm-space of theories.

Definition (of the apm-space of finitely axiomatizable theories)

Let **Fin** be the set of all finitely axiomatizable theories. The subspace  $\langle Fin, \sqsubseteq, d, D \rangle$  of the space  $\langle Th, \sqsubseteq, d, D \rangle$  is the apm-space of all finitely axiomatizable theories.

#### Definition (of the apm-space of complete theories)

Let **Comp** be the class of all complete theories. The subspace  $\langle$ **Comp**,  $\sqsubseteq$ , *d*, *D* $\rangle$  of the space  $\langle$ **Th**,  $\sqsubseteq$ , *d*, *D* $\rangle$  is the apm-space of all complete theories.

#### Fact

In (**Comp** $, \sqsubseteq, d, D)$  it is the case that:

- $d(T_1, T_2) = D(T_1, T_2)$ ,
- $T_1 \sqsubseteq T_2$  if and only if  $T_1 = T_2$ ,
- **Comp** is a metric space.

#### Theorem

In the space of all theories  $\langle \mathbf{Th}, \sqsubseteq, d, D \rangle$  for every  $T \in Th$ :

$$T = \bigcup PR(T) \,.$$

Moreover:

$$\begin{split} d(\mathrm{T}_1,\mathrm{T}_2) &= 0 & \Longleftrightarrow \ \mathrm{T}_1 \bigcirc \mathrm{T}_2 \\ & \Longleftrightarrow \ \exists_{\mathrm{T} \in \mathrm{PR}(\mathsf{Th})} \mathrm{T} \in \mathrm{T}_1 \cap \mathrm{T}_2 \\ & \longleftrightarrow \ \mathrm{cons}(\mathrm{T}_1 \uplus \mathrm{T}_2) \,. \end{split}$$

### Definition (of tangent theories) Theories $T_1$ and $T_2$ are tangent iff

```
|T_1 \cap T_2 \cap PR(\textbf{Th})| = 1 \, .
```

#### Theorem

Theories  $T_1$  and  $T_2$  are tangent iff  $cons(T_1 \uplus T_2)$  and  $T_1 \uplus T_2 \in \textbf{Comp}$ .

### Definition (of an essentially incomplete theory)

A theory T is essentially incomplete iff it is incomplete and every its finite extension is incomplete.

### Theorem

The following statements are equivalent:

- ► T is essentially incomplete,
- T is tangent to no element of **Fin**,
- ▶  $PR(Th) \cap T \neq \emptyset$  and  $PR(Fin) \cap T = \emptyset$ .

### True sentences

### Definition (of true sentences of a given theory)

- ► Suppose that T is a first-order theory which was built to axiomatize some distinguished (standard) structure 𝔐.
- ► By the set of all true sentences we will mean the of all these sentences of *L* which are satisfied by M:

$$\mathbf{V} \coloneqq \{F \in \operatorname{Sent}(\mathcal{L}) \mid \mathfrak{M} \models F\}.$$
(11)

Verisimilitude as an approximate distance

#### Fact

For any theory T completeness of V entails:

 $D(\mathbf{T}, \mathbf{V}) = d(\mathbf{T}, \mathbf{V}) + \operatorname{diam}(\mathbf{T}).$ 

# Verisimilitude as an approximate distance

- Verisimilitude is to reflect the distance of a theory from truth.
- Thus, by the previous fact, it hinges upon d(T, V) and diam(T)
- ► Verisimilitude attains maximal value (say 1) when T = V, that is in case d(T, V) = 0 and diam(T) = 0.

#### Definition (of a verisimilitude measure)

A verisimilitude measure is a function Vs:  $\mathbf{Th} \longrightarrow [0, 1]$  such that:

$$Vs(T) \coloneqq 1 - (t \cdot d(T, \mathbf{V}) + i \cdot diam(T)).$$
 (df Vs)

where  $0 < i \le t \le 1$  correspond to the weight assigned to the closeness of T from truth and the degree of incompleteness of the theory.

The End of Part VI